



## AN EXACT PENALTY APPROACH FOR MATHEMATICAL PROGRAM WITH EQUILIBRIUM CONSTRAINTS.

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**Abstract:** This paper presents an exact penalty approach to solve the mathematical problems with equilibrium constraints (MPECs). This work is based on the smoothing functions introduced in [3] but it does not need any complicated updating rule for the smoothing penalty parameters. Some numerical academic experiments are carried out to show the efficiency and robustness of this new approach. Two generic applications are also considered: the binary quadratic programs and simple number partitioning problems.

**Keywords:** Optimization; nonlinear programming; exact penalty function.



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Mathematical programs with equilibrium constraints (MPECs) represent an optimization problem including a set of parametric variational inequality or complementary constraints. In this paper, we consider optimization problems with complementary constraints, in the following form

$$(P) \begin{cases} f^* = \min f(x, y) \\ \langle x, y \rangle = 0 \\ (x, y) \in D \end{cases} \quad (1)$$

where  $f : \mathfrak{R}^{2n} \rightarrow \mathfrak{R}$  is continuously differentiable and  $D = [0, v]^{2n}$ .  $\langle \cdot, \cdot \rangle$  denotes the inner product on  $\mathfrak{R}^n$ .

We made this choice for  $D$  only to simplify the exposition, one can consider any bounded set  $D$ .

Due to the presence of complementary constraints, there is no feasible point satisfying all inequality constraints strictly so that the usual nonlinear programming constraint qualification such as Mangasarian-Fromovitz constraint qualification (MFCQ) is violated at any feasible point.

Many smoothing and relaxation methods for solving (P) have already been proposed in the literature [2, 5, 7, 8, 10]. In this study, we propose a smoothing technique to regularize the complementary constraints based on [3], we replace each constraint

$$x_i y_i = 0$$

by

$$\theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) \leq 1$$

where the parameterized function  $\theta_\varepsilon : \mathfrak{R}_+ \rightarrow [0, 1]$  is at least  $C^2$  and satisfies :

$$\theta_\varepsilon(x) \begin{cases} \approx 1 & \text{if } x \neq 0 \\ = 0 & \text{if } x = 0 \end{cases}$$

Then we define a penalty scheme to solve the problem. To avoid the updating parameter problem, we will consider  $\varepsilon$  as some new optimization variable.

This paper is organized as follows. In section 2, we present some preliminaries and assumptions on the smoothing functions and introduce the penalty method. We prove under some mild assumptions and exact penalty property in section 3 and present numerical experiments concerning academic MPECs of small sizes in section 4. The last section presents a large set of numerical experiments considering binary quadratic programs and simple number partitioning problems.

## 2 Preliminaries

In this section, we present some preliminaries concerning the regularization and approximation process. We consider functions  $\theta_\varepsilon (\varepsilon > 0)$  with the following properties:

1.  $\theta_\varepsilon$  is nondecreasing, strictly concave and continuously differentiable,
2.  $\forall \varepsilon > 0, \theta_\varepsilon(0) = 0,$
3.  $\forall x > 0, \lim_{\varepsilon \rightarrow 0} \theta_\varepsilon(x) = 1,$
4.  $\lim_{\varepsilon \rightarrow 0} \theta'_\varepsilon(0) > 0,$
5.  $\exists m > 0, \exists \varepsilon_0 > 0 \forall x \in [0, v], \forall \varepsilon \in ]0, \varepsilon_0], |\partial_\varepsilon \theta_\varepsilon(x)| \leq \frac{m}{\varepsilon^2}$

For  $\varepsilon = 0$ , we set  $\theta_0(0) = 0$  and  $\theta_0(x) = 1, \forall x \neq 0$ . Examples of such functions are:



$$\begin{aligned}
 (\theta_\varepsilon^1): \quad \theta_\varepsilon(x) &= \frac{x}{x + \varepsilon} \\
 (\theta_\varepsilon^{w_1}): \quad \theta_\varepsilon(x) &= (1 - e^{-\frac{x}{\varepsilon}})^k, \text{ for } k \leq 1 \\
 (\theta_\varepsilon^{\log}): \quad \theta_\varepsilon(x) &= \frac{\log(1+x)}{\log(1+x+\varepsilon)}
 \end{aligned}$$

Using function  $\theta_\varepsilon$ , we obtain the relaxed following problem :

$$(P_\varepsilon) \quad \begin{cases} f^* = \min f(x, y) \\ \theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) \leq 1, \quad i = 1, \dots, n \\ (x, y) \in D \end{cases} \tag{2}$$

**Remark 2.1**  $\langle x, y \rangle = 0 \Rightarrow \forall \varepsilon > 0, \theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) \leq 1$ . Thus any feasible point for  $(P_\varepsilon)$  is also feasible for  $(P)$  and then  $\forall \varepsilon > 0, f_\varepsilon^* \leq f^*$ .

We first transform the inequality constraints into equality constraints, by introducing some slacks variables  $e_i$  :

$$\theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) + e_i - 1 = 0, \quad e_i \geq 0 \quad i = 1, \dots, n. \tag{3}$$

The problem  $(P_\varepsilon)$  becomes:

$$(\tilde{P}_\varepsilon) \quad \begin{cases} \min f(x, y) \\ \theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) + e_i - 1 = 0, \quad i = 1, \dots, n \\ (x, y, e) \in D \times [0, 1]^n \end{cases} \tag{4}$$

Indeed each  $e_i$  can not exceed 1.

**Remark 2.2** The limit problem  $(\tilde{P}_\varepsilon)$  for  $\varepsilon = 0$

$$(\tilde{P}_\varepsilon) \quad \begin{cases} \min f(x, y) \\ \theta_\varepsilon(x_i) + \theta_\varepsilon(y_i) + e_i - 1 = 0, \quad i = 1, \dots, n \\ e_i \in [0, 1], \quad i = 1, \dots, n \end{cases} \tag{5}$$

is equivalent to  $(P)$ .

Until now, this relaxation process was introduced in [3]. To avoid the updating of parameters problem, we define the penalty functions  $f_\sigma$  on  $D \times [0, 1] \times [0, \bar{\varepsilon}]$ :

$$f_\sigma(x, y, e, \varepsilon) = \begin{cases} f(x, y) & \text{if } \varepsilon = \Delta(x, y, e, \varepsilon) = 0, \\ f(x, y) + \frac{1}{2\varepsilon} \Delta(x, y, e, \varepsilon) + \sigma\beta(\varepsilon) & \text{if } \varepsilon > 0, \\ +\infty & \text{if } \varepsilon = 0 \text{ and } \Delta(x, y, e, \varepsilon) \neq 0 \end{cases}$$

where  $\Delta$  measures the feasibility violation  $\Delta(z, \varepsilon) = \|G_\varepsilon(z)\|^2$  where  $(G_\varepsilon(z))_i = (\theta_\varepsilon(x) + \theta_\varepsilon(y) + e - 1)_i$  and



$$z = (x, y, e).$$

The function  $\beta : [0, \bar{\varepsilon}] \rightarrow [0, \infty)$  is continuously differentiable on  $(0, \bar{\varepsilon}]$  with  $\beta(0) = 0$ .

**Remark 2.3**  $\forall z \in D', \Delta(z, 0) = 0 \Leftrightarrow z$  feasible for  $\tilde{P} \Leftrightarrow (x, y)$  feasible for  $(P)$ .

Then we consider the following problem:

$$(P_\sigma) \quad \begin{cases} \min f_\sigma(x, y, e, \varepsilon) \\ (x, y, e, \varepsilon) \in D \times [0, 1]^n \times [0, \bar{\varepsilon}] \end{cases} \quad (6)$$

From now on, we will denote

$$D' = D \times [0, 1]^n \quad (7)$$

**Definition 2.1** We say that the Mangasarian-Fromovitz condition [9] for  $P_\sigma$  holds at  $z \in D'$  if  $G'_\varepsilon(z)$  has full rank and there exists a vector  $p \in \mathfrak{R}^n$  such that  $G'_\varepsilon(z)p = 0$  and

$$p_i \quad \begin{cases} > 0 & \text{if } z_i = 0 \\ < 0 & \text{if } z_i = w_i \end{cases}$$

with

$$w_i = \begin{cases} v & \text{if } i \in \{1 \dots 2n\} \\ 1 & \text{if } i \in \{2n+1 \dots 3n\} \end{cases}$$

**Remark 2.4** This regularity condition can be replaced by one of those proposed in [11].

### 3 The smoothing technique

The following theorem yields a condition to find a solution for  $(P_\sigma)$ . It also proves a direct link to  $(P)$ :

**Theorem 3.1** We suppose that  $z \in D'$  satisfies the Mangasarian-Fromovitz condition, and that

$$\beta'(\varepsilon) \geq \beta_1 > 0 \text{ for } 0 < \varepsilon < \bar{\varepsilon}.$$

- i) If  $\sigma$  is sufficiently large, there is no KKT point of  $P_\sigma$  with  $\varepsilon > 0$ .
- ii) For  $\sigma$  sufficiently large, every local minimizer  $(z^*, \varepsilon^*), (z^* = (x^*, y^*, e^*))$  of the problem  $(P_\sigma)$  has the form  $(z^*, 0)$ , where  $(x^*, y^*)$  is a local minimizer of the problem  $(P)$ .

**Proof:**

i) Let  $(z, \varepsilon)$  be a Kuhn Tucker point of  $P_\sigma$ , then there exist  $\lambda$  and  $\mu \in \mathfrak{R}^{3n+1}$  such that:

$$\begin{aligned} (i) \quad & \nabla l(z, \varepsilon) = \nabla f_\sigma(z, \varepsilon) + \lambda - \mu = 0, \\ (ii) \quad & \min(\lambda, z_i) = \min(\mu, w_i - z_i) = 0, \quad i = 1 \dots 3n \\ (iii) \quad & \mu_{3n+1} = \min(\lambda_{3n+1}, \bar{\varepsilon} - \varepsilon) = 0, \end{aligned} \quad (8)$$

where  $\nabla f_\sigma$  is the gradient of  $f_\sigma$  with respect to  $(z, \varepsilon)$ .

Assume that there exists a sequence of KKT points  $(z_k, \varepsilon_k)$  of  $P_{\sigma_k}$  with  $\varepsilon_k \neq 0, \forall k$  and  $\lim_{k \rightarrow +\infty} \sigma_k = +\infty$ . Since  $D'$  is bounded and closed, up to a subsequence, we have



$$\lim_{k \rightarrow +\infty} \varepsilon_k = \varepsilon^*$$

$$\lim_{k \rightarrow +\infty} z_k = z^*$$

(8.i) yields to  $\partial_\varepsilon f_{\sigma_k}(z_k, \varepsilon_k) + \lambda_{3n+1} - \mu_{3n+1} = 0$ . So that  $\partial_\varepsilon f_{\sigma_k}(z_k, \varepsilon_k) \leq 0$ .

Then, if we denote  $\Delta_k = \Delta(z_k, \varepsilon_k)$ , we have

$$\begin{aligned} \partial_\varepsilon f_{\sigma_k} &= -\frac{1}{4\varepsilon_k^2} \Delta_k + \frac{1}{2\varepsilon_k} \partial_\varepsilon \Delta_k + \sigma_k \beta'(\varepsilon_k) \\ &= -\frac{1}{4\varepsilon_k^2} \Delta_k + \frac{1}{\varepsilon_k} (\theta_\varepsilon(x_k) + \theta_\varepsilon(y_k) + e_k + 1)(\partial_\varepsilon \theta_\varepsilon(x_k) + \partial_\varepsilon \theta_\varepsilon(y_k)) + \sigma_k \beta'(\varepsilon_k) \leq 0 \end{aligned}$$

Multiplying by  $4\varepsilon^3$ , we obtain

$$4\varepsilon_k^2 (\theta_\varepsilon(x_k) + \theta_\varepsilon(y_k) + e_k - 1)(\partial_\varepsilon \theta_\varepsilon(x_k) + \partial_\varepsilon \theta_\varepsilon(y_k)) + 4\varepsilon_k^3 \sigma_k \beta'(\varepsilon_k) \leq \varepsilon_k \Delta_k$$

Since  $\Delta_k, \theta_\varepsilon$  and  $\varepsilon^2 \partial_\varepsilon \theta_\varepsilon$  are bounded (by definition (v)),  $\sigma_k \rightarrow \infty$  when  $k \rightarrow \infty$ . We have  $\varepsilon^* = 0$ .

(ii) Let  $\sigma$  sufficiently large and  $(z^*, \varepsilon^*)$  a local minimizer for  $(P_\sigma)$ . If  $(z^*, \varepsilon^*)$  satisfies the Magasarian-Fromovitz condition, then  $(z^*, \varepsilon^*)$  is a Kuhn-Tucker points for  $f_\sigma$ . By (i), we conclude that  $\varepsilon^* = 0$ .

Let  $V$  be a neighborhood of  $(z^*, 0)$ , for any  $z$  feasible for  $\tilde{P}$  such that  $(z, 0) \in V$  we have

$$f_\sigma(z^*, 0) \leq f_\sigma(z, 0) = f(x, y) < +\infty \tag{9}$$

(since  $\Delta(z, 0) = 0$ ).

We can conclude that  $\Delta(z^*, 0) = 0$ , otherwise  $f_\sigma(z^*, 0)$  would be  $+\infty$ . So that  $\langle x^*, y^* \rangle = 0$  and  $(x^*, y^*)$  is a feasible point of  $(P)$ .

Back to (9)  $f(x^*, y^*) = f_\sigma(z^*, 0) \leq f_\sigma(z, 0) = f(x, y)$ . Therefore  $(x^*, y^*)$  is a local minimizer for  $(P)$ .

**Remark 3.1** The previous theorem is still valid if we consider penalty functions of the form

$$f_\sigma(x, y, e, \sigma) = f(x, y) + \alpha(\varepsilon)\Delta(x, y, e, \varepsilon) + \sigma\beta(\varepsilon) \tag{10}$$

with  $\alpha(\varepsilon) > \frac{1}{2\varepsilon}$ .

### 4 Numerical results

In this section we consider some preliminary results obtained with the approach described in the previous section. We used the SNOPT solver [6] for the solution on the AMPL optimization platform [1]. In all our tests, we use the same function  $\beta$  defined by  $\beta(\varepsilon) := \sqrt{\varepsilon}$ .

We consider various MPECs where the optimal value is know [4]. Tables 1 and 2 summarizes our different informations concerning the computational effort of the SNOPT, by using respectively  $\theta^{w1}$  and  $\theta_1$  function:

- Obj.value : is the optimal value





- it : correspond to the total number of iterations
- (Obj.) and (grad.) : correspond to the total number of objective function evaluations and objective function gradient evaluations
- (constr.) and (jac.) : give respectively the total number of constraints and constraints gradient evaluations.

<i>Problem</i>	<i>Data</i>	<i>Obj.val.</i>	<i>it</i>	<i>Obj.</i>	<i>grad</i>	<i>constr.</i>	<i>Jac</i>
<i>bard1</i>	<i>no</i>	17	331	193	192		
<i>desilva</i>	(0,0)	-1	892	655	656	655	656
	(2,2)	-1	448	416	415	416	415
<i>Df1</i>	<i>no</i>	$3.26e-12$	3	28	27	28	27
<i>Bilevel1</i>	(25,25)	5	470	214	213		
	(50,50)	5	295	168	169		
<i>Bilevel2</i>	(0,0,0,0)	-6600	232	55	54		
	(0,5,0,20)	-6600	180	56	55		
	(5,0,15,10)	-6600	-6599.9	331	97	96	
<i>flp4</i>	<i>flp4-1.dat</i>	$1.9e-29$	66	9	8		
	<i>flp4-2.dat</i>	$3.08e-29$	66	9	8		
	<i>flp4-3.dat</i>	$1.1e-28$	66	9	8		
<i>gauvin</i>	<i>no</i>	0	184	71	70		
<i>jr1</i>	<i>no</i>	0.5	1175	814	813		
<i>scholtes1</i>	1	2	12	10	9	10	9
<i>hs044</i>	<i>no</i>	14.97	375	101	100		
<i>nash1</i>	(0,0)	0	0	2	1		
	(5,5)	$1.72e-17$	13	13	12		
	(10,10)	$2.24e-12$	12	12	11		
	(10,0)	$4.29e-12$	12	11	10		
	(0,10)	$1.46e-13$	13	12	11		
<i>qpec1</i>	<i>no</i>	80	1249	443	442		
<i>liswet1-inv</i>	<i>liswet1-050</i>	0.0929361	215	126	125		
<i>Stack1</i>	0	-3266.67	27	26			
	100	-3266.67	7	17	16		
	200	-3266.67	7	17	16		
<i>Water-net</i>	<i>Water-net.dat</i>	931.1	2070	886	885	886	885

Table 1: using the  $\theta^{w_1}$  function



Problem	Data	Obj.val.	it	Obj.	grad	constr.	Jac
<i>bard1</i>	<i>no</i>	17	433	248	247		
<i>desilva</i>	(0,0)	-1	7	255	254	255	254
<i>Df1</i>	<i>no</i>	0	657	961	960	961	960
<i>gauvin</i>	<i>no</i>	9.5e-05	164	82	81	82	81
<i>Bilevel1</i>	(25,25)	5	401	190	198		
	(50,50)	5	391	183	182		
<i>Bilevel2</i>	(0,0,0,0)	-6600	2458	487	486	487	486
	(0,5,0,20)	-6600	2391	727	721		
	(5,0,15,10)	-6600	2391	727	721		
<i>hs044</i>	<i>no</i>	17.08	617	261	260	261	260
<i>jr1</i>	<i>no</i>	0.5	67	54	53		
<i>nash1</i>	(0,0)	3.35e-13	203	111	110		
	(5,5)	6.7e-24	146	71	70		
	(10,10)	2.3e-17	133	85	84		
	(10,0)	8.1e-16	379	238	237		
	(0,10)	2.37e-18	1228	848	847		
<i>qpecl</i>	<i>no</i>	80	1895	518	517		
<i>liswet1-inv</i>	<i>liswet1-050</i>	0.028	3559	462	461		
<i>scholtes1</i>	1	2	51	106	105	106	105
<i>Stack1</i>	0	-3266.67	64	58	57		
	100	-3266.67	30	32	31		
	200	-3266.67	30	32	31		
<i>Water-net</i>	<i>Water-net.dat</i>	931.369	919	282	281	282	281

Table 2: using the  $\theta^1$  function

We remark that by considering  $\theta^{w1}$  or  $\theta^1$  we obtain the optimal know value in almost all the considered test problems.

## 5 Application to simple partitioning problem and binary quadratic problems

In this section, we consider two real applications : the simple number partitioning and binary quadratic problems. These two classes of problems are know to be NP hard. We use here our approach as a simple heuristic to obtain local solutions.

### 5.1 Application to simple partitioning problem

The number partitioning problem can be stated as a quadratic binary problem. We model this problem as follows.

We consider a set of numbers  $S = \{s_1, s_2, s_3, \dots, s_m\}$ . The goal is to divide  $S$  into two subsets such that the subset sums are as close to each other as possible. Let  $x_j = 1$  if  $s_j$  is assigned to subset 1, 0 otherwise. Then  $sum_1$ , subset

1's sum, is  $sum_1 = \sum_{j=1}^m s_j x_j$  and the sum for subset 2 is  $sum_2 = \sum_{j=1}^m s_j - \sum_{j=1}^m s_j x_j$ . The difference in the sums is then

given by

$$diff = \sum_{j=1}^m s_j - 2 \sum_{j=1}^m s_j x_j = c - 2 \sum_{j=1}^m s_j x_j. \quad (c = \sum_{j=1}^m s_j).$$

We will minimize the square of this difference

$$diff^2 := (c - 2 \sum_{j=1}^m s_j x_j)^2,$$

We can rewrite  $diff^2$  as



$$\text{diff}^2 = c^2 + 4x^T Qx,$$

where

$$q_{ii} = s_i(s_i - c), \quad q_{ij} = s_i s_j.$$

Dropping the additive and multiplicative constants, our optimization problem becomes simply

$$UQP \begin{cases} \min x^T Qx \\ x \in \{0,1\}^n \end{cases}$$

We rewrite the problem as the follows:

$$UQP \begin{cases} \min x^T Qx \\ x_i(1-x_i) = 0 \end{cases}$$

With this formulation, the proposed algorithm can be used to get some local solutions for (UPQ).

We considered modest-sized random problems ( $m = 25$  and  $m = 75$ ). Five instances of each size were generated with the elements drawn randomly from the interval (50,100).

Table 3 summarizes Table the obtained results. For each instance, we used 100 different initial points generated randomly from the interval  $[0,1]$ :

- Best diff : corresponds to the best value of  $\left| \sum_{i=1}^{100} (Q * \text{round}(x_i) - 0.5 * c) \right|$
- Integrality measure : correspond to the  $\max_i | \text{round}(x_i) - x_i |$
- Nb: correspond to the number of tests such that the best sum is satisfied.
- $Nb_{10}$  : correspond to the number of tests such that the sum :  $\left| \sum_{i=1}^{100} (Q * \text{round}(x_i) - 0.5 * c) \right| \leq 10$





<i>Problem</i>	<i>Best diff</i> $(\theta^1, \theta^{w_1})$	<i>Nb</i> $(\theta^1, \theta^{w_1})$	<i>Integrality measure</i> $(\theta^1, \theta^{w_1})$	<i>Nb<sub>10</sub></i> $(\theta^1, \theta^{w_1})$
<i>NP25.1</i>	(1,0)	(1,2)	(0.011,0)	(15,15)
<i>NP25.2</i>	(0,0)	(2,2)	(0.0055,0.005)	(16,14)
<i>NP25.3</i>	(0,0)	(1,1)	(0,0)	(16,14)
<i>NP25.4</i>	(0,0)	(1,2)	(0,0)	(22,22)
<i>NP25.5</i>	(0,0)	(1,4)	(0.008,0.0045)	(11,10)
<i>NP75.1</i>	(0,0)	(1,2)	(0.003,0)	(14,14)
<i>NP75.2</i>	(0,0)	(2,1)	(0,0)	(15,15)
<i>NP75.3</i>	(0,0)	(1,1)	(0,0)	(17,17)
<i>NP75.4</i>	(0,0)	(2,2)	(0,0)	(18,18)
<i>NP75.5</i>	(0,1)	(1,1)	(0,0)	(17,17)

Table 3: Results on Partitioning Problem using the  $\theta^1$  and  $\theta^{w_1}$  function

## 5.2 Application to binary quadratic problems

We consider some test problems from the Biq Mac Library [12]. These problems are written in the simple following formulation:

$$\begin{aligned} \min y^T Q y \\ y \in \{0,1\}^n \end{aligned}$$

where  $Q$  is a symmetric matrix of order  $n$ .

We used ten instances with  $n = 100$  for the  $Q$  matrix generated using the following restrictions:

- diagonal coefficients in the range  $[-100,100]$ ,
- off-diagonal coefficients in the range  $[-50,50]$ ,
- All coefficients are integers,
- seeds  $1,2,\dots,10$ .

The third column (Nbop) of table 4 report the number of optimal realizations (times we obtained the known optimal value) with 100 different initial points generated randomly from the interval  $[0,1]$ . The fourth column precise the obtained value when it is different from the known optimal value.



<i>Problem</i>	<i>Know.value</i>	<i>Nbop</i> ( $\theta^l, \theta^{w_1}$ )	<i>Foundvalue</i> ( $\theta^l, \theta^{w_1}$ )
<i>be100.1</i>	-19412	(17,14)	
<i>be100.2</i>	-17290	(14,12)	
<i>be100.3</i>	-17565	(9,13)	
<i>be100.4</i>	-19125	(9,14)	
<i>be100.5</i>	-15868	(2,2)	
<i>be100.6</i>	-17368	(31,31)	
<i>be100.7</i>	-18629	(0,0)	(-18473,-18475)
<i>be100.8</i>	-18649	(1,1)	
<i>be100.9</i>	-13294	(0,0)	(-13248,-13248)
<i>be100.10</i>	-15352	(11,4)	

Table 4: Results on Biq Mac test problems using the  $\theta^l$  and  $\theta^{w_1}$  functions

Using  $\theta^{w_1}$  or  $\theta^l$  we obtain the optimal know value for almost all the instances. We obtain a local solutions for only two examples. For each instance, the algorithm found a solution and needs  $< 1s$  for the resolution.

## 6 conclusion

In this paper, we introduced an exact penalty approach to solve the mathematical program with equilibrium constraints. We proved a convergence result under suitable constraint qualification conditions and performed some numerical experiments. We initially tested our approach on some tests from the library MaMPEC. Then, we considered some examples from the Biq Mac Library and some randomly generated partitioning problems. We used two different smoothing functions and our limited numerical tests gave very promising results (almost the same result for each one).

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