# **Traveling Salesman Problem: A Case Study**

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#### **ABSTRACT**

In this paper assignment based integer linear formulation presented for solving traveling salesman problem. Unfortunately, the assignment model can lead to infeasible solutions. Infeasibility removes by introducing additional constraints. Then this linear problem solved by open source software.

#### **General Terms**

Optimization technique, Linear Programming, Assignment problem. TSP, combinatorial optimization problem

#### Keywords

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# 1. INTRODUCTION

The Traveling Salesman Problem (TSP) is a classical combinatorial optimization problem, which is simple to state but very difficult to solve. The problem is to find the shortest possible tour through a set of N vertices so that each vertex is visited exactly once. This problem is known to be NPcomplete, and cannot be solved exactly in polynomial time because the number of possible routes increases factorally with the number of cities. In other words, for 5 cities there are 5! (120) possible routes and for 6 cities there are 6! (720). A salesman who had to visit the 15 European capitals in this example would have 1,307,674,368,000 (over 1.3 trillion) possible routes to consider. This traditional approach become impractical in terms of computer memory and speed constraints. So to solve this problem different researchers use heuristic, metaheuristic and optimal method like dynamic programming, linear programming. In this paper traveling salesman problem solved like assignment problem using linear programming approach. The constraints require that the salesman must enter and leave each city exactly once. Unfortunately, the assignment model can lead to infeasible solutions. Consequently, additional constraints must be included in order to eliminate subtour solutions. There are a number of ways to accomplish this.

#### 2. REVIEW OF LITERATURE

The Traveling Salesman Problem (TSP) is one of the important subjects which have been widely addressed extensively by mathematicians and computer scientists [1]. Many approaches have been tried to solve TSP problem. Among them, there are polynomial-sized linear programming formulation [2], Genetic Algorithm [3,4,5,6], a mixed integer linear programming formulation and dynamic programming [7,8,9], NP–complete problem[10,11], Neural Network [12]. The multiple traveling salesman problem (mTSP) is a generalization of the well-known traveling salesman problem (TSP), where more than one salesman can be used in the solution [13]. Historically, researchers have suggested a

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multitude of heuristic algorithms, including genetic algorithms (GA's) [14] for solving TSP [15]. Recently studied approaches to solve TSP includes Knowledge based multiple inversion and knowledge based neighborhood swapping GA algorithms [16], hybridized GA for integrated TSP and quadratic assignment problem [17], two-level GA for clustered TSP and large scale TSP [18], parallel GA program implementation on multicomputer cluster [19].

# 3. ASSIGNMENT BASED FORMULATION

Starting from his home, a salesman wishes to visit each of (n-1) other cities and return home at minimal cost. He must visit each city exactly once and it costs cij to travel from city i to city j.

We may be tempted to formulate his problem as the assignment problem:

$$x_{ij} = \begin{cases} 1 & \text{if he goes from city } i \text{ to city } j, \\ 0 & \text{otherwise,} \end{cases}$$

$$\mathbf{Minimize} \sum_{i=1}^{n} \sum_{j=1}^{n} c_{ij} \ X_{ij}$$

Subject to

$$\sum_{i=1}^{n} x_{ij} = 1 \quad (j = 1, 2, \dots, n),$$

$$\sum_{j=1}^{n} x_{ij} = 1 \quad (i = 1, 2, \dots, n),$$

Additional constraints

$$x_{i,j} \ge 0$$
  $(i = 1, 2, ..., n; j = 1, 2, ..., n).$ 

The constraints require that the salesman must enter and leave each city exactly once. Unfortunately, the assignment model can lead to infeasible solutions. Infeasibility removes by additional constraints discussed in next section.

#### 4. EXPERIMENT SETUP

To explain the procedure to solve traveling salesman problem we consider the following example in which we have five cities 1, 2, 3, 4 and 5.

From	1	2	3	4	5
1	-	3	6	2	3

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2	3	-	5	2	3
2 3 4 5	3 6 2	5	-	6	4
4	2	2	6	-	6
5	$\begin{bmatrix} 2 \\ 3 \end{bmatrix}$	3	4	6	-

In the above problem we have 5 cities, let us designate the visit of  $i^{th}$  city to  $j^{th}$  city by the decision variable  $x_{ij}$ , i=1,2,3,4,5 and j=1,2,3,4,5 but  $i\neq j$ .

#### **Mathematical Formulation**

#### MIN Cost=

```
3*x_{12}+6*x_{13}+2*x_{14}+3*x_{15}+3*x_{21}+5*x_{23}+2*x_{24}+3*x_{25}+6*x_{31}+6*x_{32}+6*x_{34}+4*x_{35}+2*x_{41}+2*x_{42}+6*x_{43}+6*x_{45}+3*x_{51}+3*x_{52}+4*x_{53}+6*x_{54};
```

#### Subject to Constraints:

```
 \begin{aligned} &(\mathbf{x}_{12} + \mathbf{x}_{13} + \mathbf{x}_{14} + \mathbf{x}_{15}) = 1 \\ &(\mathbf{x}_{21} + \mathbf{x}_{23} + \mathbf{x}_{24} + \mathbf{x}_{25}) = 1 \\ &(\mathbf{x}_{31} + \mathbf{x}_{32} + \mathbf{x}_{34} + \mathbf{x}_{35}) = 1 \\ &(\mathbf{x}_{41} + \mathbf{x}_{42} + \mathbf{x}_{43} + \mathbf{x}_{45}) = 1 \\ &(\mathbf{x}_{51} + \mathbf{x}_{52} + \mathbf{x}_{53} + \mathbf{x}_{54}) = 1 \\ &(\mathbf{x}_{21} + \mathbf{x}_{31} + \mathbf{x}_{41} + \mathbf{x}_{51}) = 1 \\ &(\mathbf{x}_{12} + \mathbf{x}_{32} + \mathbf{x}_{42} + \mathbf{x}_{52}) = 1 \\ &(\mathbf{x}_{13} + \mathbf{x}_{23} + \mathbf{x}_{43} + \mathbf{x}_{53}) = 1 \\ &(\mathbf{x}_{14} + \mathbf{x}_{24} + \mathbf{x}_{34} + \mathbf{x}_{54}) = 1 \\ &(\mathbf{x}_{15} + \mathbf{x}_{25} + \mathbf{x}_{35} + \mathbf{x}_{45}) = 1 \\ &\forall \ x_{ij} \ge 0 \ (i = 1, 2, ..., 5; j = 1, 2, ..., 5). \end{aligned}
```

To solve the ILPP, Software LINGO Version 11 was used. Figure 1 and Figure 2 illustrate the snap shots of construction of ILPP and generated solution by software LINGO 11, respectively.

### **Solution:**

```
\begin{array}{l} x_{12}\!\!=\!x_{24}\!\!=\!x_{41}\!\!=\!x_{35}\!\!=\!x_{53}\!\!=\!1 \\ x_{13}\!\!=\!x_{14}\!\!=\!x_{15}\!\!=\!x_{21}\!\!=\!x_{23}\!\!=\!x_{25}\!\!=\!x_{31}\!\!=\!\!x_{32}\!\!=\!x_{34}\!\!=\!x_{42}\!\!=\!x_{43}\!\!=\!x_{45}\!\!=\!x_{51}\!\!=\!x_{52}\!\!=\!x_{54}\!\!=\!0 \end{array}
```

It is clear from the produced result assignment solution to route the salesman through disjoint subtours of the cities(1-2-4-1 and 3-5-3) instead of on a single trip or tour (figure 3)

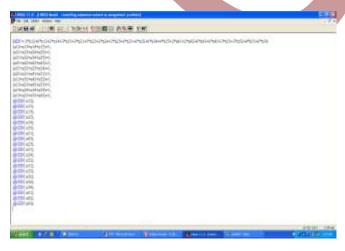


Fig 1: Formulation (Lingo 11 Software)

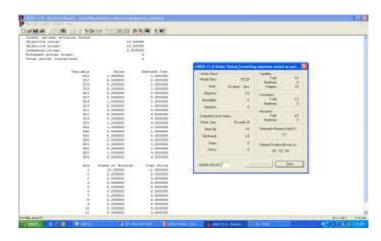


Fig 2: Output Generated Using Lingo 11 Software

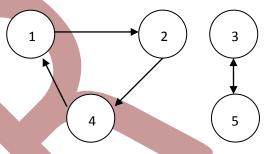


Fig 3: disjoint subtours

Unfortunately, the assignment model can lead to infeasible solutions. Infeasibility removes by additional constraints like  $x_{13}+x_{15}+x_{23}+x_{25}+x_{43}+x_{45}>=1;.....(1)$ 

This inequality ensures that at least one leg of the tour connects cities 1, 2, and 4 with cities 3 and 5. In general, if a constraint of this form is included for each way in which the cities can be divided into two groups, then subtours will be eliminated. Now again above ILPP solved with additional constraint (1) using Software LINGO Version 11 that produce 1 tour cycle (fig 4).

#### **Solution:**

Minimum Cost z=16  $x_{14}=x_{23}=x_{35}=x_{42}=x_{51}=1$ 

 $x_{12} = x_{13} = x_{15} = x_{21} = x_{24} = x_{25} = x_{31} = x_{32} = x_{34} = x_{41} = x$ 

 $x_{43} = x_{45} = x_{52} = x_{53} = x_{54} = 0$ 

Above result produce 1 tour cycle 1-4-2-3-5-1.

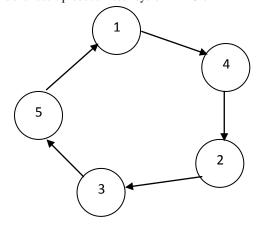


Fig 4: One Tour Cycle

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#### 5. CONCLUSION

In this paper assignment based integer linear formulation presented for solving traveling salesman problem. Unfortunately, the assignment model can lead to infeasible solutions. Infeasibility remove by introducing additional constraints. Then this linear problem solved by open source software. Due to the constraint of length of the paper only one case have been reported here. However, more than ten problems have been solved using this approach.

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